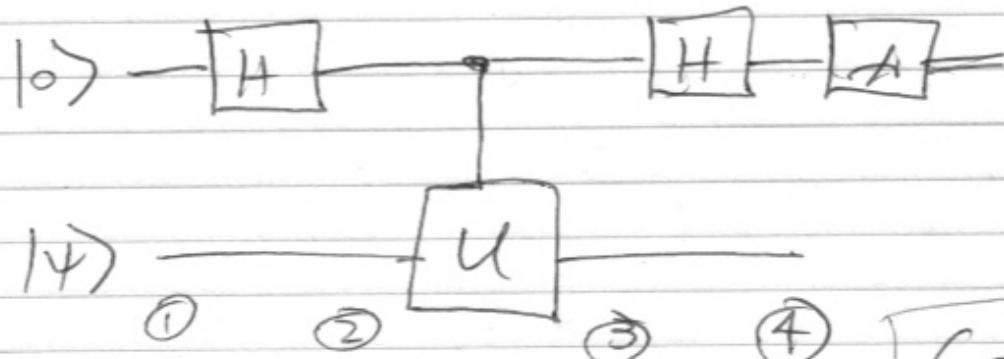


Lecture 7

①

We can apply reasoning from before to a generic controlled-unitary.

Consider the following circuit



Analyze again step by step

$$\begin{aligned} C-U = \\ |0\rangle &\otimes I + \\ |1\rangle &\otimes U \end{aligned}$$

① $|0\rangle|1\rangle$

\rightarrow ② $|+\rangle|1\rangle = \frac{|0\rangle+|1\rangle}{\sqrt{2}}|1\rangle = \frac{|0\rangle|1\rangle+|1\rangle|1\rangle}{\sqrt{2}}$

\rightarrow ③ $\frac{|0\rangle|1\rangle+|1\rangle U|1\rangle}{\sqrt{2}}$

(2)

$$\rightarrow \textcircled{4} \quad \frac{|+\rangle|+\rangle + |-\rangle|-\rangle}{\sqrt{2}}$$

probability to get outcome zero:

$$p(0) = \left\| \left(|0\rangle \langle 0| \otimes I \right) \frac{1}{\sqrt{2}} (|+\rangle|+\rangle + |-\rangle|-\rangle) \right\|_2^2$$

$$= \left\| |0\rangle \left(\frac{|+\rangle + |-\rangle}{2} \right) \right\|_2^2$$

$$= \left\| |0\rangle \left(\frac{I+U}{2} \right) |+\rangle \right\|_2^2$$

$$= \left[(0| \langle +| \left(\frac{I+U}{2} \right) \right] \left[|0\rangle \left(\frac{I+U}{2} \right) |+\rangle \right]$$

$$= \frac{1}{4} \langle +| (I+U^+) (I+U) |+\rangle$$

$$= \frac{1}{4} \langle +| 2I + U + U^+ |+\rangle$$

$$= \frac{1}{4} (2 + \langle +|U|+\rangle + \langle +|U^+|+\rangle)$$

$$= \frac{1}{2} (1 + \operatorname{RE}[\langle +|U|+\rangle])$$

(3)

Check: Is this a legitimate probability?

Consider that

$$\begin{aligned} \text{RE}\{\langle +|U|\psi \rangle\} &\leq |\langle +|U|\psi \rangle| \\ &\leq \cancel{\sqrt{\|\psi\|_2 \|\psi\|_2}} = 1 \end{aligned}$$

~~or $\text{RE}\{\langle +|U|\psi \rangle\}$~~

Same reasoning gives

$$-\text{RE}\{\langle +|U|\psi \rangle\} \leq 1$$

$$\Rightarrow \text{RE}\{\langle +|U|\psi \rangle\} \geq -1$$

$$\Rightarrow \frac{1}{2}(1 + \text{RE}\{\langle +|U|\psi \rangle\}) \in [0, 1] \quad \checkmark$$

(4)

What is the probability of getting outcome 1?

$$\begin{aligned}
 p(1) &= \|((1\rangle\langle 1|01)\frac{1}{\sqrt{2}}(1\rangle|1\rangle + 1\rangle|0\rangle))\|_2^2 \\
 &= \|(1\rangle\left(\frac{I-U}{2}\right)|1\rangle)\|_2^2 \\
 &= \frac{1}{2}(1 - \text{RE}[\langle 1|U|1\rangle])
 \end{aligned}$$

Question: Describe a q-algorithm for estimating $\text{RE}[\langle 1|U|1\rangle]$ w/r addtive error ϵ w/ success probability at least $1-\delta$.

(5)

Now let us consider a special input state & a special choice of unitary

pick input state to be

$$|\psi\rangle = |\psi\rangle \otimes |\phi\rangle \quad +$$

unitary to be $U = \text{SWAP}$, which acts as

$$\text{SWAP } |\psi\rangle \otimes |\phi\rangle = |\phi\rangle \otimes |\psi\rangle \\ (\text{see homework})$$

then

$$\begin{aligned} \langle \psi | U | \psi \rangle &= \langle \psi | \otimes \langle \phi | \text{SWAP} |\psi\rangle \otimes |\phi\rangle \\ &= (\langle \psi | \otimes \langle \phi |) (|\phi\rangle \otimes |\psi\rangle) \\ &= \langle \psi | \phi \rangle \langle \phi | \psi \rangle = |\langle \psi | \phi \rangle|^2 \end{aligned}$$

This is called the SWAP test)

Implies that this algorithm
can estimate the inner product
of $|q\rangle$ & $|f\rangle$.

One of the most important
tasks in q. machine learning.

Idea is to encode
classical information into
high-dimensional states +
then use this basic
algorithm to estimate
the inner product $|\langle q | f \rangle|^2$.

If the states are complex (i.e.,
difficult to simulate on a classical
computer), then the idea is
that this could present

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an advantage for QML.

In particular, the method of support vector machines evaluates inner products of high-dimensional vectors as a basic subroutine.

There is another way to estimate inner products & it is called the destructive SWAP test.

Let us consider the case of two qubits $|1\rangle$ & $|+\rangle$

consider from homework 2
that

$$\begin{aligned}
 & \text{Tr}[\text{SWAP} |\psi\rangle\langle e| \otimes |\phi\rangle\langle \phi|] \\
 &= \text{Tr}[|\psi\rangle\langle e| |\phi\rangle\langle \phi|] \\
 &= \langle \phi | \psi \rangle \langle e | \phi \rangle \\
 &= |\langle \phi | \psi \rangle|^2
 \end{aligned}$$

we can write first line as

↗ $\langle \phi | \text{SWAP} | \psi \rangle$

where $|\psi\rangle = |e\rangle \otimes |\phi\rangle$

this is an expectation of
the SWAP observable.

From homework, we know that
for qubits,

$$\text{SWAP} = \Phi^+ + \Phi^- + \Psi^+ - \Psi^-$$

where we use the

shorthand $\Phi^+ \equiv |\Phi^+\rangle\langle\Phi^+|$,
etc.

$$\Rightarrow \langle 4 | \text{SWAP} | 4 \rangle$$

$$= \langle 4 | \Phi^+ | 4 \rangle + \langle 4 | \Phi^- | 4 \rangle$$

$$+ \langle 4 | \Psi^+ | 4 \rangle - \langle 4 | \Psi^- | 4 \rangle$$

$$= |\langle \Phi^+ | 4 \rangle|^2 + |\langle \Phi^- | 4 \rangle|^2$$

$$+ |\langle \Psi^+ | 4 \rangle|^2 - |\langle \Psi^- | 4 \rangle|^2$$

Given this, what is a
q. algorithm to estimate

$$\langle 4 | \text{SWAP} | 4 \rangle ?$$

(back to first week of
class)

Set $i = 1$.
Answer: Perform Bell measurement,
set $\Xi_i = 1$ if outcome
 $\Phi^+, \Psi^-, \Psi^+ + \Xi_i = -1$ if
outcome Ψ^-

~~Step~~ Increment i .

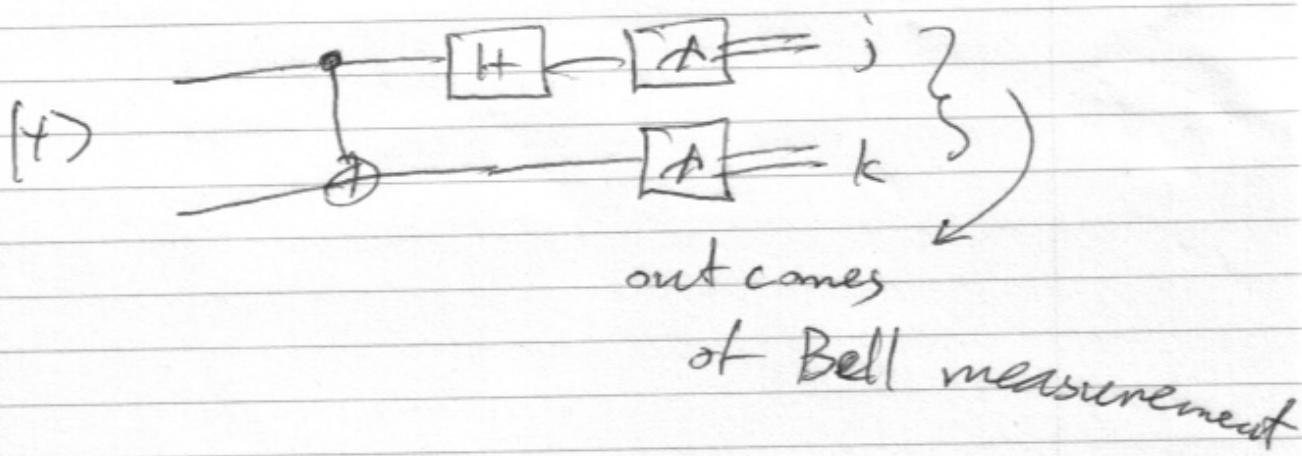
(10)

repeat for sufficiently many
times as required by
Hoeffding bound.

Set $\bar{Z}^n = \frac{1}{n} \sum_{i=1}^n Z_i$

as estimate of $\langle f | \text{SWAP} | f \rangle$

Circuit looks like this:



called destructive SWAP test

because state is destroyed in
the process.

Multi-qubit generalization available
on page 9 of arXiv:2309.02515