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Lecture 22

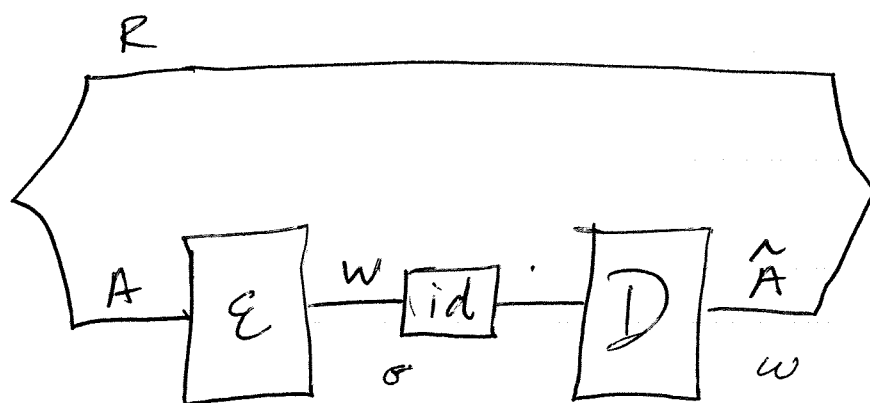
quantum data compression

(AKA Schumacher compression)

one-shot version of a q. data compression protocol.

Given is a density operator ρ_A w/ purification ψ_{RA}^P .

General protocol is as follows:



where E is an encoding or compression channel & D is a decoding or decompression channel

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Protocol is an $(\frac{|w|}{\epsilon}, \epsilon)$ protocol

if
$$F(\psi_{RA}^p, (\underbrace{D_{W \rightarrow \hat{A}}}_{(*)} \circ \epsilon_{A \rightarrow W}^{\psi_{RA}^p})) \geq 1 - \epsilon$$

one-shot ϵ -data compression
limit

$$QDL^\epsilon(p) = \inf_{(\epsilon, D), |w|} \left\{ \log_2 |w| : (*) \text{ holds?} \right\}$$

Asymptotic quantity

$$QDL(p) = \sup_{\epsilon \in (0,1)} \limsup_{n \rightarrow \infty} \frac{QDL^\epsilon(p^{\otimes n})}{n}$$

Strong converse quantity

$$\widetilde{QDL}(p) = \inf_{\epsilon \in (0,1)} \liminf_{n \rightarrow \infty} \frac{QDL^\epsilon(p^{\otimes n})}{n}$$

We always have

$$QDL(p) \geq \widetilde{QDL}(p)$$

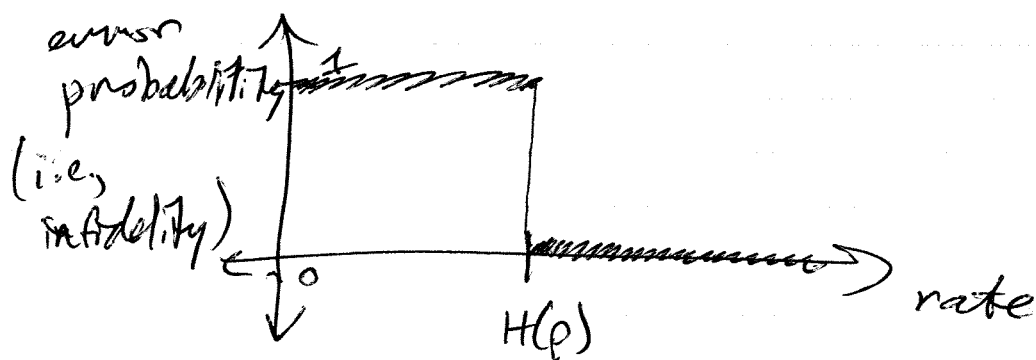
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Schumacher Data compression theorem

$$QDL(p) = \widetilde{QDL}(p) = H(p)$$

(strong converse due to Winter)

visualization of this theorem in $n \rightarrow \infty$ limit



above data compression limit,

there exists a sequence of protocols w/ error prob. tending to zero.

below QDL, every scheme has error probability tending to one.

∴ QDL is a sharp dividing line between possible & impossible.

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Let us start w/ strong converse.

Key tool is Rényi mutual information

$$\tilde{I}_\alpha(C; D)_\rho = \inf_{\sigma_D} \tilde{D}_\alpha(\rho_{CD} \| \rho_C \otimes \sigma_D)$$

~~Property~~ where

$$\tilde{D}_\alpha(\omega \| \tau) = \frac{1}{\alpha-1} \log_2 \text{Tr} \left[\left(\tau^{\frac{1-\alpha}{2\alpha}} \omega \tau^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right]$$

Key properties

$$\tilde{I}_\alpha(C; D)_\rho \leq 2 \log_2 |D| \quad \begin{matrix} \leftarrow \text{for} \\ \alpha \in [1/2, 1) \\ \cup (1, \infty) \end{matrix}$$

Data processing under local channels $\downarrow$

$$\tilde{I}_\alpha(C; D)_\rho \geq \tilde{I}_\alpha(C'; D')_\omega$$

where

$$\omega_{C'D'} = (N_{C \rightarrow C'} \otimes M_{D \rightarrow D'}) (\rho_{CD})$$

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pseudo-continuity:

$$\tilde{I}_\beta(C;D)_\rho - \tilde{I}_\alpha(C;D)_\sigma$$

$$\geq \frac{\alpha}{1-\alpha} \log F(\rho_{CD}, \sigma_{CD})$$

$$\text{for } \alpha \in (1/2, 1) \text{ \& } \beta = \frac{\alpha}{2\alpha-1}$$

Proofs: see arXiv:1506.02635

Also:

$$\tilde{I}_\alpha(C;D)_\psi = 2H \frac{1}{2\alpha-1}(C)_\psi$$

for pure state ψ_{CD} \&
 $\alpha \in (0, 1) \cup (1, \infty)$

in PACT book

$$\alpha \in (1/2, 1) \quad \beta = \frac{\alpha}{2\alpha-1} \quad (6)$$

Now use all these properties

to prove a one-shot bound;

for an arbitrary (W, ϵ) protocol

$$2 \log_2 |W| \geq \tilde{I}_\beta(R; W)_\sigma \quad (\text{dim. bound})$$

$$\geq \tilde{I}_\beta(R; \hat{A})_w \quad (\text{data processing})$$

$$\geq \tilde{I}_\alpha(R; \hat{A})_{\varphi^p} \quad (\text{pseudo-continuity})$$

$$+ \frac{\alpha}{1-\alpha} \log_2 F(\varphi_{R\hat{A}}^p, w_{R\hat{A}})$$

$$\geq \tilde{I}_\alpha(R; \hat{A})_{\varphi^p}$$

$$+ \frac{\alpha}{1-\alpha} \log_2 (1-\epsilon)$$

$$= 2 H_{\frac{1}{2\alpha-1}}(A)_\rho + \frac{\alpha}{1-\alpha} \log_2 (1-\epsilon)$$

$$\Rightarrow \log_2 |W| \geq H_{\frac{1}{2\alpha-1}}(A)_\rho + \frac{\alpha}{2(1-\alpha)} \log_2 (1-\epsilon)$$

holds for arbitrary protocol

$\Rightarrow$ one-shot bound

$$QDL^\epsilon(\rho) \geq H_{\frac{1}{2\alpha-1}}(A)_\rho + \frac{\alpha}{2(1-\alpha)} \log_2 (1-\epsilon)$$

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IID case:

$$\frac{1}{n} QDL^\varepsilon(p^{\otimes n})$$

$$\geq \frac{1}{n} H_{\frac{1}{2\alpha-1}}(p^{\otimes n}) + \frac{\alpha}{2n(1-\alpha)} \log_2(1-\varepsilon)$$

$$= H_{\frac{1}{2\alpha-1}}(p) + \frac{\alpha}{2n(1-\alpha)} \log_2(1-\varepsilon)$$

Take $n \rightarrow \infty$

$$\liminf_{n \rightarrow \infty} \frac{1}{n} QDL^\varepsilon(p^{\otimes n}) \geq H_{\frac{1}{2\alpha-1}}(p)$$

holds $\forall \alpha \in (1/2, 1)$

take $\alpha \rightarrow 1$ limit

$$\Rightarrow \liminf_{n \rightarrow \infty} \frac{1}{n} QDL^\varepsilon(p^{\otimes n}) \geq H(p)$$

$$\Rightarrow \widetilde{QDL}(p) \geq H(p)$$

↑
von Neumann
entropy