

# Lecture 16

①

## Fast Fourier Transform

Recall that the DFT is

$$c_k = \sum_{n=0}^{N-1} y_n \exp(-i \frac{2\pi kn}{N})$$

Python program that we gave before has a for-loop for each coefficient &  $N$  terms in the sum, implying  $N^2$  ops to get all coefficients.

If we are not willing to wait for more than 1 billion ops, then we can do a DFT for

$$N^2 = 10^9 \Rightarrow N \approx 32000 \text{ samples}$$

This is not very much just about one second of audio. The FFT is a much better way for doing this, is covered by

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easiest to describe when # of samples is a power of two.

Let  $N = 2^m$  w/  $m$  an integer

Consider the sum in the DFT equation & divide it into 2 groups, even & odd terms

Consider even terms first, where

$n = 2r$  w/  $r \in \{0, \dots, \frac{1}{2}N - 1\}$

then 
$$E_k = \sum_{r=0}^{\frac{1}{2}N-1} y_{2r} \exp\left(-i \frac{2\pi k(2r)}{N}\right)$$

call this  $\rightarrow = \sum_{r=0}^{\frac{1}{2}N-1} y_{2r} \exp\left(-i \frac{2\pi kr}{N/2}\right)$

$E_{k \bmod N/2}$  This is actually just another DFT w/  $N/2$  samples instead of  $N$

look @ odd terms now:

$$\sum_{r=0}^{\frac{1}{2}N-1} y_{2r+1} \exp\left(-i \frac{2\pi k(2r+1)}{N}\right) = \rightarrow$$

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$$= e^{-i2\pi k/N} \sum_{r=0}^{N/2-1} y_{2r+1} \exp\left(-i \frac{2\pi k r}{N/2}\right)$$

$$= e^{-i2\pi k/N} O_k$$

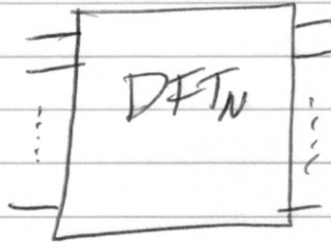
where  $O_k$  is another DFT w/  $N/2$  samples  
 call this  $O_{k \bmod N/2}$   
 So this means that

$$C_k = \sum_{k \bmod N/2} e^{-i2\pi k/N} O_{(k \bmod N/2)}$$

$$= \sum_{k \bmod N/2} W_N^k O_{k \bmod N/2}$$

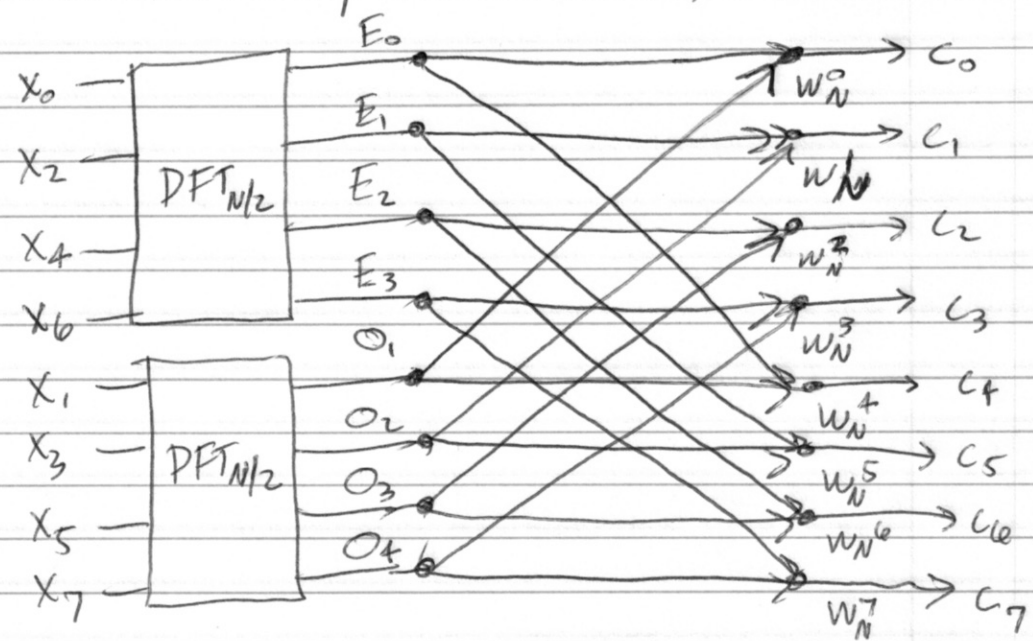
We can write what we have written in terms of a diagram.

Let



denote an  $N$ -point DFT. The outputs are the DFT of the input

Then to compute  $C_0$ , we do



Thus the famous "butterfly" diagram for the FFT.

It specifies the recursion & the starting point is  $DFT_4$  which is just the identity transformation.

At the ~~nth~~ ~~stage~~  $m$ th stage there are  $O(N)$  calculations.

There are  $\log N$  stages, meaning that the ~~total~~ FFT requires  $O(N \log N)$  operations. Near linear!

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Fig. if we have  $N =$  1 million samples to process, then naive way requires

$N^2 = 10^{12}$  operations, which is nontrivial for a typical computer. & not practical

But  $N \log N = 10^7$  which is reasonable (under one second)

Inverse DFT can be done in the same way.

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FFT in Python provided by module `numpy.fft`

bring up `1-fft.py`

`rfft` computes FFT for a set of real numbers, only returns the 1st half because the other half are complex conjugates by symmetry



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can also calculate Fourier transforms  
of complex data w/ fft & ifft

& 2D transforms w/ rfft2, irfft2,  
fft2, ifft2

also functions for higher  
dimensions

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Chapter C - Solving ordinary differential  
equations

1st-order differential equations

Consider 
$$\frac{dx}{dt} = \frac{2x}{t} + \frac{3x^2}{t^3}$$

It is not separable & is also nonlinear,  
so we need to resort to the computer  
to solve it.

general form for a 1st-order  
differential equation  $\exists \frac{dx}{dt} = f(x,t)$

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to calculate a full solution, we need a boundary condition, e.g., the value of  $x$  @ one particular value of  $t$  (usually  $t=0$ ).

### Euler's method

Suppose we have  $\frac{dx}{dt} = f(x,t)$   
to solve

& an initial condition.

Then Taylor expand  $x(t+h)$  about  $t$

$$\begin{aligned}x(t+h) &= x(t) + h \frac{dx}{dt} + \frac{h^2}{2} \frac{d^2x}{dt^2} + \dots \\ &= x(t) + h f(x,t) + o(h^2)\end{aligned}$$

If we neglect  $o(h^2)$  terms

then we get

$$x(t+h) = x(t) + h f(x,t)$$

so if we know  $x$  @ time  $t$ , we

can just use this equation to iterate  
if  $h$  is small enough, then this does pretty well.

called Euler's method

Example  $\frac{dx}{dt} = -x^3 + \sin t$

Bring up 2-euler.py

No one really uses Euler's method because there is a more accurate approach that requires very little extra effort to program. These are the Runge-Kutta methods.

The error in one step of the Euler method is to leading order

$$\frac{1}{2}h^2 \frac{d^2x}{dt^2}$$

Error for many steps ~~is then~~ from  $t=a$  to  $t=b$  using step size  $h$  is then

$N = \frac{b-a}{h}$  values of  $t$  for steps are @  $t_k = a + kh$



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$x_k$  is corresponding value for  $x$

Then the total error is

$$\sum_{k=0}^{N-1} \frac{1}{2} h^2 \left( \frac{d^2 x}{dt^2} \right)_{\substack{x=x_k, \\ t=t_k}} = \frac{1}{2} h \left[ \sum_{k=0}^{N-1} h \left( \frac{df}{dt} \right)_{\substack{x=x_k, \\ t=t_k}} \right]$$

$$\approx \frac{1}{2} h \int_a^b \frac{df}{dt} dt$$

$$\approx \frac{1}{2} h [f(x(b), b) - f(x(a), a)]$$

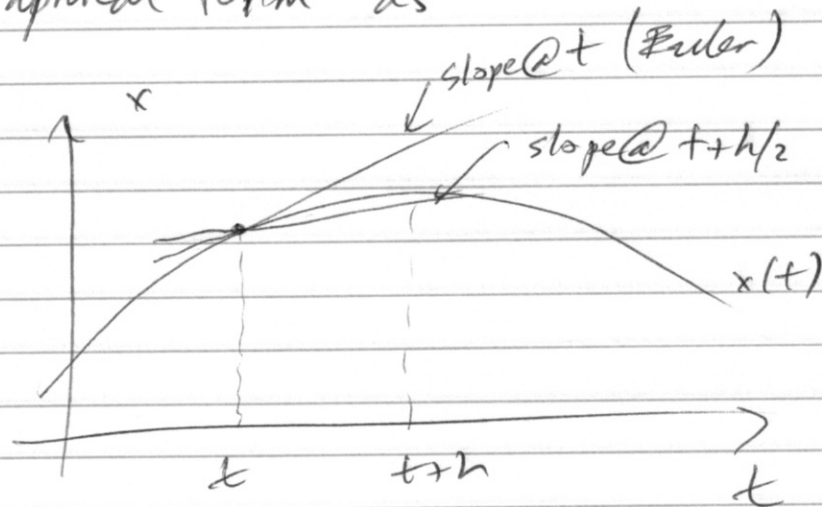
approx. is good if  $h$  is small.

error is linear in  $h$ , so that

if we make it smaller, then error is smaller.

Runge-Kutta method can do much better. (these are really a set of methods)

Can express Euler's method in graphical form as



curve is  $x(t)$ , diff. eq.  ~~$\frac{dx}{dt} = f(x,t)$~~

$\frac{dx}{dt} = f(x,t)$  says  
that slope of solution is  
equal to  $f(x,t)$ .

So Euler method extrapolates slope to  
time  $t+h$ , giving an estimate of  
 $x(t+h)$

2nd order Runge-Kutta method is to use the slope  
@ midpoint  $t+h/2$  to extrapolate to  
 $t+h$

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Perform a Taylor expansion of  $x(t+h)$   
about  $t + \frac{1}{2}h$ :

$$x(t+h) = x(t + \frac{1}{2}h) + \frac{1}{2}h \left( \frac{dx}{dt} \right)_{t + \frac{1}{2}h} + \frac{1}{8}h^2 \left( \frac{d^2x}{dt^2} \right)_{t + \frac{1}{2}h} + o(h^3)$$

$x(t)$  has expansion

$$x(t) = x(t + \frac{1}{2}h) - \frac{1}{2}h \left( \frac{dx}{dt} \right)_{t + \frac{1}{2}h} + \frac{1}{8}h^2 \left( \frac{d^2x}{dt^2} \right)_{t + \frac{1}{2}h} + o(h^3)$$

subtract 2nd from 1st to get

$$\begin{aligned} x(t+h) &= x(t) + h \left( \frac{dx}{dt} \right)_{t + \frac{1}{2}h} + o(h^3) \\ &= x(t) + h f(x(t + \frac{1}{2}h), t + \frac{1}{2}h) + o(h^3) \end{aligned}$$

$\Rightarrow$  that the order  $h^2$  term has disappeared! & error term is now order  $h^3$

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Problem is that this approach requires an estimate of  $x(t + \frac{1}{2}h)$  which we don't have (we only know  $x(t)$ ). Instead approximate  $x(t + \frac{1}{2}h)$  ~~using~~ using Euler's method:

$$x(t + \frac{1}{2}h) = x(t) + \frac{1}{2}h f(x, t)$$

if then substitute. So we get

$$k_1 = h f(x, t)$$

$$k_2 = h f(x + \frac{1}{2}k_1, t + \frac{1}{2}h)$$

$$x(t+h) = x(t) + k_2$$

Final Runge-Kutta estimate

This is a "second order" method accurate to order  $h^2$  w/ error of order  $h^3$ .

How do we know this when we approximated  $x(t + \frac{1}{2}h)$  w/ Euler's method?

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Consider that (T.E.  $f(x + \frac{1}{2}k_1, t + \frac{1}{2}h)$   
around  $x(t + \frac{1}{2}h)$ )

$$f(x(t) + \frac{1}{2}k_1, t + \frac{1}{2}h)$$

$$= f(x(t + \frac{1}{2}h), t + \frac{1}{2}h)$$

$$+ [x(t) + \frac{1}{2}k_1 - x(t + \frac{1}{2}h)] \left( \frac{\partial f}{\partial x} \right)_{x(t + \frac{1}{2}h), t + \frac{1}{2}h}$$

$$+ O([x(t) + \frac{1}{2}k_1 - x(t + \frac{1}{2}h)]^2)$$

From before we have

$$x(t + \frac{1}{2}h) = x(t) + \frac{1}{2}h f(x, t) + O(h^2)$$

$$= x(t) + \frac{1}{2}k_1 + O(h^2)$$

$$\Rightarrow x(t) + \frac{1}{2}k_1 - x(t + \frac{1}{2}h) = O(h^2) +$$

$$f(x(t) + \frac{1}{2}k_1, t + \frac{1}{2}h)$$

$$= f(x(t + \frac{1}{2}h), t + \frac{1}{2}h) + O(h^2)$$

$$\Rightarrow k_2 = h f(x(t + \frac{1}{2}h), t + \frac{1}{2}h) + O(h^3)$$

So error introduced by Euler approximation is order  $h^3$



Bring up 3-rk2.py

run for different #s of  
points -  $N = 10, 20, 50, 100$

Can continue the Runge-Kutta  
method further. Take more  
Taylor expansions to cancel out

higher order terms. Popular method is

fourth order Runge-Kutta method  
which is still relatively simple  
to program:

$$k_1 = h f(x, t)$$

$$k_2 = h f(x + \frac{1}{2}k_1, t + \frac{1}{2}h)$$

$$k_3 = h f(x + \frac{1}{2}k_2, t + \frac{1}{2}h)$$

$$k_4 = h f(x + k_3, t + h)$$

$$x(t+h) \equiv x(t) + \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4)$$

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- accurate to terms of order

$h^4$   $f$  has error of order  $h^5$

- derivation of this is complicated  
but the final equations are simple

- just 5 equations but 3 orders of  
 $h$  more accurate than Euler's method

- 4th order R-K is the method  
of choice for solving ODEs on  
computer.

Bring up 4-rk4.py

run w/  $N = 10, 20, 50, 100$

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## Finding solutions over an infinite range

some times want to go out to  $t = \infty$   
idea is to  
use a change of variables as  
before.

$$\text{Let } u = \frac{t}{1+t} \quad \text{or} \quad t = \frac{u}{1-u}$$

so that  $u \rightarrow 1$  as  $t \rightarrow \infty$

re write diff. eq. w/ chain rule as

$$\frac{dx}{du} \frac{du}{dt} = f(x, t) \quad \text{or}$$

$$\frac{dx}{du} = \frac{dt}{du} f\left(x, \frac{u}{1-u}\right)$$

$$\text{But } \frac{dt}{du} = \frac{1}{(1-u)^2} f\left(x, \frac{u}{1-u}\right)$$

$$\text{so } \frac{dx}{du} = \frac{1}{(1-u)^2} f\left(x, \frac{u}{1-u}\right)$$

Let  $g(x,u) = (1-u)^2 f(x, \frac{u}{1-u})$

so  $\frac{dx}{du} = g(x,u)$  if then just use techniques from before

Example: Suppose we want to solve

$\frac{dx}{dt} = \frac{1}{x^2+t^2}$  from  $t=0$  to  $t=\infty$   
w/  $x=1$  @  $t=0$

then instead solve

$g(x,u) = \frac{1}{(1-u)^2} \frac{1}{x^2 + \frac{u^2}{(1-u)^2}}$   
 $= \frac{1}{x^2(1-u)^2 + u^2}$

So solve

$\frac{dx}{du} = \frac{1}{x^2(1-u)^2 + u^2}$  from  $u=0$  to  $u=1$

Bring up 5-ode intopy