

First discuss Bell's theorem + then
No-cloning theorem.

We will give a simple more modern approach to Bell's theorem called the "CHSH game".

main reference is ~~book~~

"The Device Independent Outlook on Quantum Physics"

Valerio Scarani 1303, 3081

- Bell's theorem is considered to be one of the greatest results in quantum mechanics + perhaps one of the first results that follows the spirit of quantum information (i.e., showing a separation ^{between} the classical & quantum theories of information)

- So many in QIT "claim" how to be one of us

- shows how entanglement is different from ~~any classical notion~~

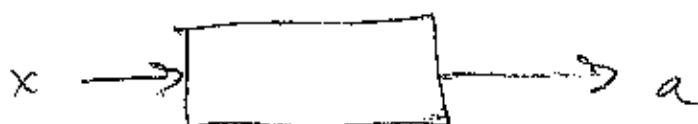
Bell's theorem also has practical consequences for secure communication using quantum resources.

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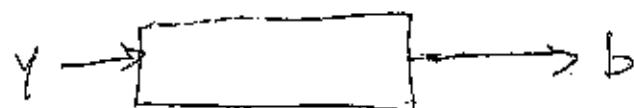
The setup is

$$x, y, a, b \in \{0, 1\}$$

Alice



Bob



spatially separated but allowed to meet before x & y are coins x & y are tossed, Alice given uniformly @ random responds w/ a & Bob w/ b .

They win if $x \wedge y = a \oplus b$

We will prove that the maximum probability of winning w/ a classical strategy is $\leq 3/4$

Quantumly, one can ~~not~~ win w/ probability $= \cos^2(\pi/8) \approx 0.85$

(3)

Proof that winning probability classically is
 $\leq \frac{3}{4}$

In order to bound this, we need to ~~not~~ consider the distribution

$$P(a, b | x, y)$$

We will use a parameter δ to describe one's "favorite explanation" for this distribution. So by expanding w/ the law of total probability, we can write

$$P(a, b | x, y) = \int d\delta \rho(\delta | x, y) P(a, b | x, y, \delta)$$

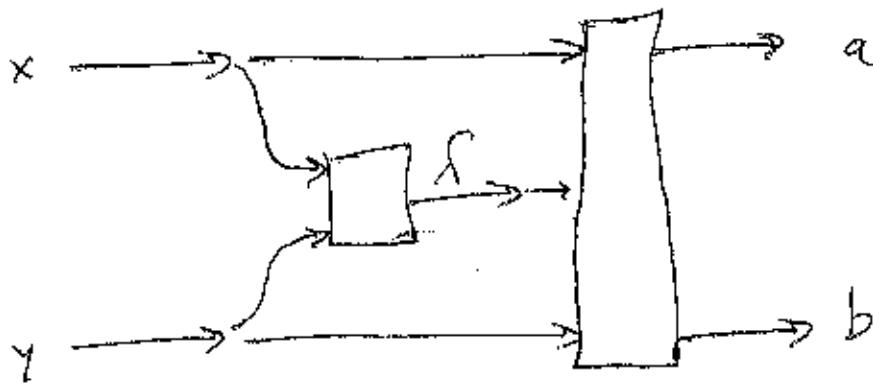
where $\rho(\delta | x, y) \geq 0$ &

$$\int d\delta \rho(\delta | x, y) = 1$$

& $P(a, b | x, y, \delta)$ are all valid distributions

~~Sketch~~ For such a general model, we can draw a picture to explain what's going on

④



If quantum theory is your favorite explanation, then

$$\rho(\lambda|x,y) = \delta(\lambda - \gamma) +$$

there are measurements

$$\{\Pi_a^{(x)}\}_a + \{\Pi_b^{(y)}\}_b \text{ such that}$$

$$P(a,b|x,y,\lambda) = \langle \Psi | \Pi_a^{(x)} \otimes \Pi_b^{(y)} | \Psi \rangle$$

i.e., redraw picture as

$$P(a,b|x,y) = \langle \Psi | \Pi_a^{(x)} \otimes \Pi_b^{(y)} | \Psi \rangle$$

Idea of Bell's theorem is to put quantum theory to the test by not taking it as the favorite & seeking out alternatives

For example, we can try out

classical explanations + see if they
can explain the results of experiments

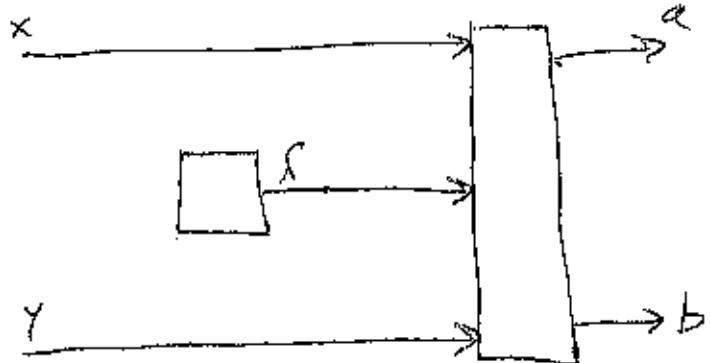
So we first need to formalize what
we mean by a classical model

- we call such a model pre-established agreement or also "local hidden variable theory" (which sounds somewhat mysterious)
general

Returning to the picture given before,
there are a few things about it
that are not consistent w/ our
understanding of how the game works

First, the parameter δ represents
correlations shared between Alice &
Bob before the game begins. So
it is not sensible that δ depends on
 x & y . So we modify the picture
 δ to be as

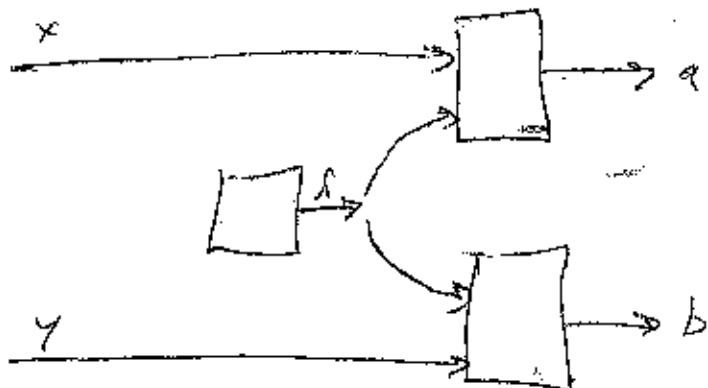
⑥



(called "measurement independence")

$$\text{i.e., } p(l|x,y) \\ = p(l)$$

Next, as part of the game, we said that Alice + Bob are spatially separated, so that ~~there should be~~ they should be acting locally on their inputs. The picture now changes as



$$\text{i.e., } P(a,b|x,y,l) = P(a|x,l) P(b|y,l)$$

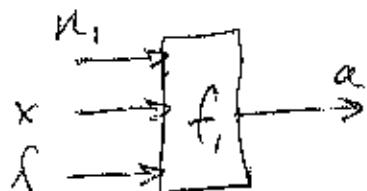
so we have now argued that any classical explanation for what's going on should have the form

$$P(a,b|x,y) = \int d\lambda p(\lambda) P(a|x,\lambda) P(b|y,\lambda)$$

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we now argue that it suffices to consider deterministic strategies for winning the CHSH game.

The maps $P(a|x,s)$ & $P(b|y,s)$ are stochastic, but we can simulate these as

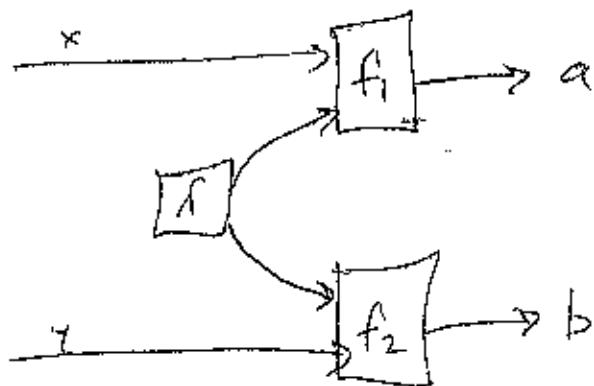


where n_i is a local noise variable independent of $x \neq 1$ & f_i is a deterministic function, so that

$$P(a|x,s) = \int_{\Omega} p(n_i) f(a|n_i, x, s) dn_i$$

same thing for $p(b|y,s)$ w/ a noise variable

But then we can ~~subsume~~ ^{n_2 & function f_2} subsume the local noise variables into s , so that the picture becomes



Furthermore, the winning probability $\overleftarrow{\text{average}}$ is a convex combination of the probabilities that a collection of deterministic strategies wins. Since an average of a set of numbers cannot be larger than the maximum of those numbers, ~~there~~ there always ~~is~~ a deterministic strategy that does just as well as a probabilistic one.

- So it suffices to obtain an upper bound on the winning probability of any deterministic strategy.

recall that they win if $x \wedge y = a \oplus b$

consider that any deterministic strategy

has $x \rightarrow a_x$

$y \rightarrow b_y$

So consider that

Is it possible to
win all of the time?
Would imply a
contradiction

x	y	$x \wedge y$	a_x	b_y	$a_x \oplus b_y$
0	0	0			$a_0 \oplus b_0$
0	1	0			$a_0 \oplus b_1$
1	0	0			$a_1 \oplus b_0$
1	1	1			$a_1 \oplus b_1$
					0

can only win $\frac{3}{4}$ of the time. ⑨

So how is it possible to win 85% of the time in QM?

First, we allow them to share the maximally entangled state

$$|\Psi\rangle_{AB} = \frac{1}{\sqrt{2}} (|0\rangle_A \otimes |0\rangle_B + |1\rangle_A \otimes |1\rangle_B)$$

Alice measures the observables

$$A_0 = \sigma_z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad \text{or} \quad A_1 = \sigma_x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

depending on x or
Bob measures

$$B_0 = \frac{\sigma_x + \sigma_z}{\sqrt{2}} \quad \text{or} \quad B_1 = \frac{\sigma_z - \sigma_x}{\sqrt{2}}$$

depending on y

How to measure observables?

can write $\sigma_z = |0\rangle\langle 0| - |1\rangle\langle 1|$

outcomes are $+1$ & -1 w/ projectors $\{|0\rangle\langle 0|, |1\rangle\langle 1|\}$

(10)

$$\text{For } \sigma_x = |+\rangle\langle+| - |- \rangle\langle-|$$

$$\text{where } |+\rangle \equiv \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$$

$$|-\rangle \equiv \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

outcomes are $+1$ & -1 and projectors are

$$\{|+\rangle\langle+, |-\rangle\langle-\}\}$$

similar kind of thing for Bob

identify a' as the measurement outcome so that

$$a' = (-1)^a \quad a' \in \{+1, -1\} \\ a \in \{0, 1\}$$

Consider that

$$\langle +|_{AB} (A_x \otimes B_y) |+\rangle_{AB}$$

is the expected value of the product of their measurement outcomes

When ~~$x, y \in \{0, 1, 10\}$~~ $x, y \in \{0, 1, 10\}$

they should report the same outcome.

so in these cases,

$\langle +|_{AB} (A_x \otimes B_y) |+\rangle_{AB} \Rightarrow$ the prob. of winning minus
prob. of losing

When ~~$x^2 + y^2 = 1$~~ , they should report different outcomes so that winning prob. - losing prob is equal to

$$- \langle + |_{AB} A_1 \otimes B_1 | + \rangle_{AB}$$

After averaging over choices of x & y , winning prob - losing prob

$$P_W - P_L =$$

$$\frac{1}{4} \langle + | A_0 \otimes B_0 + A_+ \otimes B_+ + A_- \otimes B_- \\ - A_1 \otimes B_1 | + \rangle$$

one can check that all four terms are equal to $\frac{1}{\sqrt{2}}$ \Rightarrow

$$P_W - P_L = \frac{1}{\sqrt{2}}$$

combined w/ $P_W + P_L = 1$, we get

$$P_W = \frac{1}{2} + \frac{1}{2\sqrt{2}} = \cos^2(\pi/8)$$

(12)

Tsirelson's bound

one cannot beat $\cos^2(\pi/8)$ in the CHSH game

Consider any observables A_0, A_1, B_0, B_1

w/ eigenvalues ~~$\{-1, +1\}$~~

~~$\{-1, +1\}$~~ & any state $|+\rangle$

such that $A_i^2 = B_j^2 = I$

then

$$\begin{aligned} & [A_0 \otimes B_0 + A_0 \otimes B_1 + A_1 \otimes B_0 - A_1 \otimes B_1] \\ &= 4I - [A_0, A_1][B_0, B_1] \end{aligned}$$

then use that

$$\|[A_0, A_1]\|_{\infty} \leq 2\|A_0\|_{\infty}\|A_1\|_{\infty} \leq 2$$

$$\Rightarrow |\langle + | \dots | + \rangle| \leq \sqrt{8} = 2\sqrt{2}$$

\Rightarrow bound on winning prob.